

16 - Implementations

10 stycznia 2011
12:30

We need three elements:

- 1) source of qubits (photons)
- 2) channels (free space or optical fibers)
- 3) detectors (we need single photon detectors)

Ideally we would like to prepare a single photon in a definite polarization state, send it through channel where it experiences no (or controllable) polarization transformation, and have an ideal detector which clicks if a photon enters it. In reality we do not have any of these elements:

[16.1 Source of qubits]

Single photon sources are impractical so much more common is the use of weak coherent laser pulses
Repetition rate \approx GHz

1a) - coherent state photon statistics $p_m = \frac{\mu^m}{m!} e^{-\mu} = \frac{\mu^m}{m!} e^{-\mu}$

So apart from single photons we have multiphoton contribution.

Photon number splitting attack (PNS)

- E performs photon number non-destructive measurement
- If $n=1$ she performs a "single photon" attack (she may also block this photon and send vacuum to B)

If $n > 1$ she takes away one photon and sends the remaining ones to B. After A & B choose their basis E measures her stored photon in the correct basis

We need to use weak pulses so that the multiphoton component is not too big ---

[16.2 Channels]

Light is attenuated $I(L) = I(0) 10^{-\alpha \cdot L}$

- free space $\alpha < 0,1$ dB/km
(780-850 nm, 1520-1600 nm)
(10 dB - ten times attenuation \approx 100 km)
nice

problems - weather, line of sight, beam broadening
 ideas maybe via a satellite,

- fibers $\alpha = 0.34 \text{ dB/km}$ (1330 nm)
 $\alpha = 0.12 \text{ dB/km}$ (1550 nm)

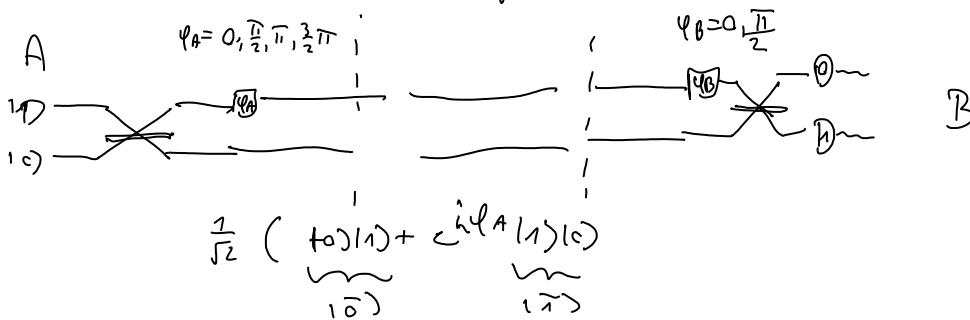
in optical regime not practical, losses too big $\sim 10 \text{ dB/km}$.

So we need infrared single photon detectors more challenging than in visible range

— fiber is birefringent, sensitive to stress and temperature variations. Polarization is subject to "random" transitions

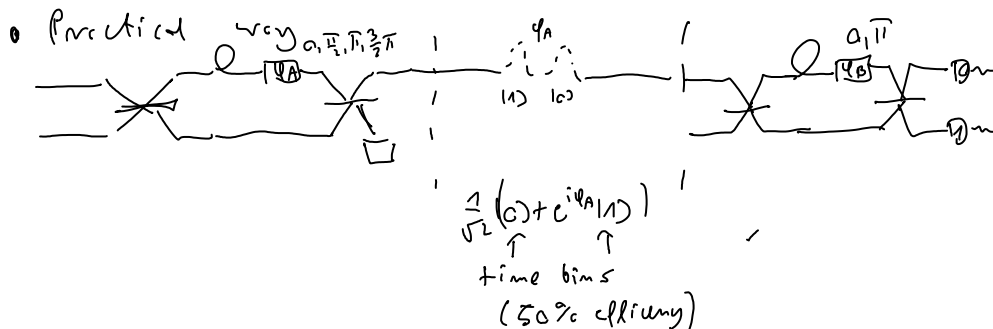
Solution: Use a different physical qubit than

- than polarization qubit



if $\phi_B = 0$ then $\phi_A = 0 \rightarrow$ perfect interference
 $\phi_A = \pi \rightarrow$
 $\phi_A = \pi/2 \rightarrow$ $p = 1/2$ $p = 1/2$
 $\phi_A = 3\pi/2 \rightarrow$ $p = 1/2$ $p = 1/2$

Exactly the same probabilities as in polarization measurement. But this requires stable interference over 100s km - not practical



At B side: $\frac{1}{\sqrt{2}} (|0\rangle + e^{i\phi_A} |1\rangle) \rightarrow \frac{1}{2} (|0, u\rangle + |0, L\rangle + e^{i\phi_A} |1, u\rangle + e^{i\phi_A} |1, L\rangle)$

$$\rightarrow \frac{1}{2} (|1, u\rangle + |0, l\rangle + e^{i\phi_A} |2, u\rangle + e^{i\phi_A} |1, l\rangle) \rightarrow$$

$$\xrightarrow{\text{PBS}} \frac{1}{2} (e^{i\phi_B} |1, u\rangle + |0, l\rangle + e^{i(\phi_A + \phi_B)} |2, u\rangle + e^{i\phi_A} |1, l\rangle)$$

Provided we observe clicks in $|1\rangle$ time bin we have interference and we receive μ -bit/s from ideal setup (overall only 25% = 50% · 50% efficiency)

[16.3 Detectors]

We need to register single photon counts

- Si APD $\eta = 50\%$ $\lambda \approx 400-1000\text{nm}$ $S \approx 100\text{Hz}$
(not suitable for infrared) $\text{count rate} \approx 15\text{MHz}$
($T = -30\text{K}$)
- InGaAs $\eta = 10\%$ $\lambda = 1000-1650\text{nm}$
count rate 0,1 MHz $S \approx 100\text{Hz}$ ($T \approx -100\text{K}$)

[16.4 Practical security analysis]

We use coherent pulses with mean intensity μ , at a repetition rate R_0 , over a fiber of length L , with attenuation coefficient α , detectors with efficiency η and dark count rate S .

At what rate K we can generate secure key?
(secure bits per second)

We attribute all errors to E and assume her technology is unlimited (lossless fibers, ideal devices etc.). We treat that errors from dark count due to E intervention...

A & B measure QBER (error per bit) and the rate R (number of clicks per second on B side)

B receives a coherent state with average number of photons

$$\mu = \mu_0 \cdot 10^{-\alpha \cdot L}$$

Detector efficiency can be regarded as additional loss

$$\mu = \mu_0 \cdot 10^{-\alpha \cdot L} \eta = \mu \cdot \tilde{\eta}$$

Probability that B registers a click from state sent

$P_{\text{transmission}} = 1 - e^{-\mu}$ by A
 \uparrow prob. of receive

But additionally B may get a click due to dark counts $P_{\text{dark}} = \frac{S}{R_0}$

So total click probability:

$P_{\text{click}} = P_{\text{trans}} + P_{\text{dark}} - \underbrace{P_{\text{trans}} P_{\text{dark}}}_{\text{usually negligible}}$

Detection rate by B:

$R = R_0 \cdot P_{\text{click}} \approx R_0(1 - e^{-\mu}) + S \approx R_0 \mu + S$

While $QBER = \frac{1}{2} \frac{S}{R} + \frac{\Sigma}{R}$
 \uparrow other sources of error (visibility of interferometer, noise in the fiber, (usually negligible for) large distances)

Information gained by E via optimal PNS attack

Let $P_{m \geq 2} = 1 - e^{-\mu_0} \approx \mu_0 e^{-\mu_0}$ be the probability of more than 1 photon sent by A

$P_{m \geq 2} \approx 1 - 1 + \mu_0 - \frac{\mu_0^2}{2} - \mu_0 + \mu_0^2 \approx \frac{\mu_0^2}{2}$

If $P_{\text{click}} \leq P_{m \geq 2}$ it means E can block single photon signals at any point multiplier signals to B or which she has full knowledge so key generation not possible (E uses ideal lossless fiber)

If $P_{\text{click}} > P_{m \geq 2}$ then E has to allow some single photon pulses to go to B.

Let f_1 be fraction of $m=1$ pulses that E transmits to B, but of course before that she performs the optimal attack (collective - be careful what collective means here ...)

Then $P_{\text{click}} = P_{m \geq 2} + f_1 P_{m=1} = \frac{R}{R_0}$

It is ... can derive f_1 :

This way we can derive f_1 :

$$f_1 = \frac{1}{\sum_{n=1}^{\infty} \binom{R}{n} p_n^{n-1} (1-p_n)} \approx \frac{\mu}{\mu_0} + \frac{\sum}{2R\mu_0} - \frac{\mu_0}{2}$$

E information per bit

$$I(A:E) = 1 - \frac{Y_{\text{multi}}}{Y_{\text{class}}} + \frac{Y_1}{Y_{\text{class}}} \cdot h(\epsilon_1)$$

↑
error rate introduced in single photon states

$$Y_1 = (1 - \frac{R_0}{R} p_{n \geq 2}) \approx 1 - \frac{\mu_0^2}{2(\mu + \frac{\sum}{R_0})}$$

- function of visibilities registered by B engineering from single photon states

$$QBER = \epsilon_1 \cdot Y_1$$

$$\epsilon_1 = \frac{QBER}{Y_1}$$

$$I(A:E) = (1 - Y_1) + Y_1 h\left(\frac{QBER}{Y_1}\right)$$

$$I(A:B) = 1 - h(QBER) = 1 - h\left(\frac{\sum}{2R}\right)$$

Key note:

$$K = \frac{1}{2} R (I(A:B) - I(A:E)) =$$

↑ because only in half of the cases basis (can be avoided by using one basis almost all of the time) will agree

$$= \frac{1}{2} R \left[(1 - h(QBER)) - (1 - Y_1) - Y_1 h\left(\frac{QBER}{Y_1}\right) \right] = \frac{1}{2} R \left[Y_1 (1 - h\left(\frac{QBER}{Y_1}\right)) - h(QBER) \right]$$

We can optimize over the choice of μ_0

Optimal $\mu_0 \propto \tilde{\gamma}$ (we need to lower intensity to lower the danger of multi-photon events)

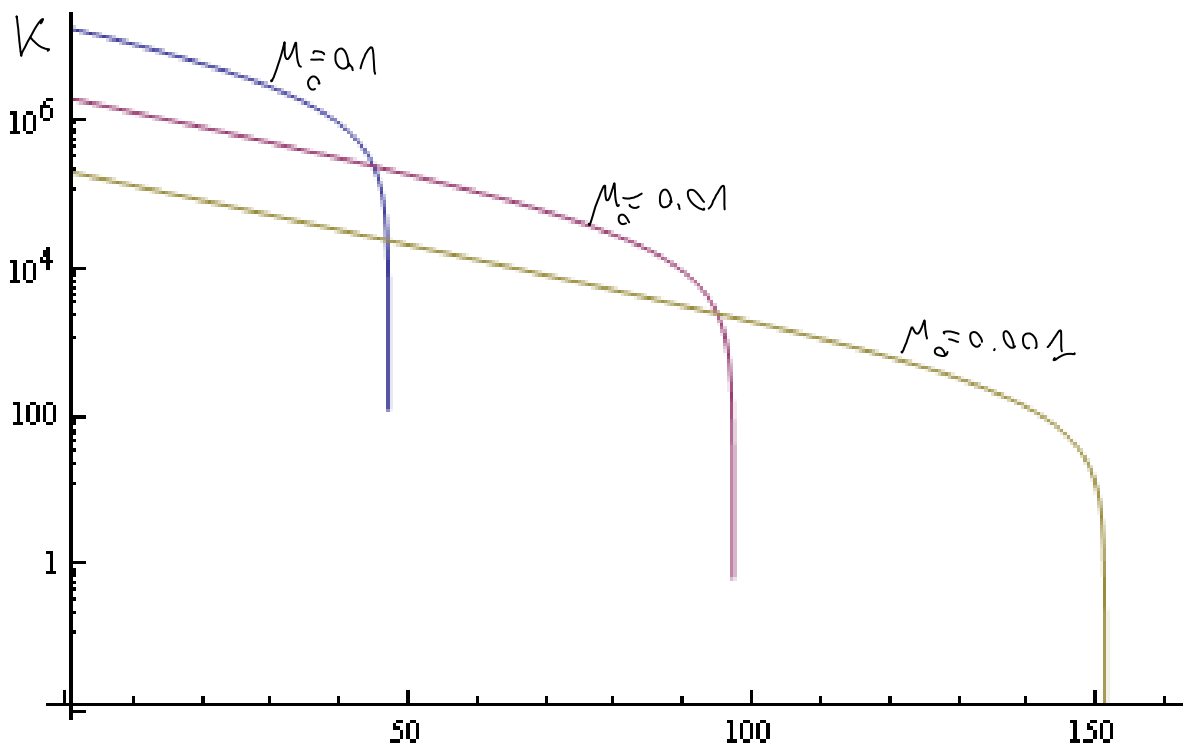
So we need to lower the intensity proportionally

to losses \downarrow hence $K \propto \tilde{\gamma}^2$

key note drops very quickly

key rate drops very quickly

A Plot: for $\alpha = 0,2 \text{ dB/km}$ $\eta = 10\%$ $S = 1 \text{ GHz}$ $\bar{R} = 16 \text{ Hz}$



[16.5 Practical (rel.)]

- Decoy states, to limit the power of PNS attacks. A sends a few different types of pulses M_1, M_2, M_3, \dots (three is enough)

This allows A & B to find exactly $\gamma_0, \gamma_1, \gamma_2, \dots$ and not only the rate R . Since PNS violates ratios between γ_i , and this ratio are ok then excludes PNS attack. We can look to

$$K \propto \eta$$

- DPSK
- Continuous variable cryptography
- MagicQ, 1d Quantum

[16.6 Quantum Hacking]

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- Imperfect devices
 - e.g. different gating times of detectors
 - Side channels - other physical channels reveal information about basis settings..
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